

Lecture 7 – Dimension Reduction and Johnson-Linderstrauss Lemma

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1 Introduction

This lecture mainly focuses on dimension reduction: Johnson-Linderstrauss Lemma and especially its distributional version. Chi-squared distribution is introduced when there is a sum of Gaussian distributed variables.

2 Last Time

Tug-of-War:

- for frequency vector $f \in \mathbb{R}$:
- pick random $\sigma_i \in \{\pm 1\}$
- $z_i = \sum_{i=1} \sigma_i f_i$
- Estimator: z^2

Tug-of-War+ : k estimators

$$z_j = \sum_{i=1} \sigma_{ij} f_i, j = 1, \dots, k$$

Estimator:

$$\frac{1}{k} \sum z_j^2$$

3 Dimension Reduction

Definition 1 (Sketching function). For $\bar{x} \in \mathbb{R}^n$, $\bar{x} = (x_1, x_2, \dots, x_n)$, a sketching function $\varphi : \mathbb{R}^n \rightarrow \mathbb{R}^k$ is defined as

$$\varphi(x) = \frac{1}{\sqrt{k}} \left(\sum \sigma_{1i} x_i, \sum \sigma_{2i} x_i, \dots, \sum \sigma_{ki} x_i \right)$$

Definition 2 (Linear Property). φ is linear if:

$$\varphi(x) + \varphi(y) = \varphi(x + y)$$

$$\varphi(x) - \varphi(y) = \varphi(x - y)$$

Estimator:

$$\varphi(x) \rightarrow \|\varphi(x)\|^2 = \frac{1}{k} \sum_{j=1}^k z_j^2$$

$$\varphi(x) = \frac{1}{\sqrt{k}}(z_1, z_2, \dots, z_k)$$

Given sketches $\varphi(x)$ and $\varphi(y)$:
we can compute

$$\|\varphi(x) - \varphi(y)\|_2^2 = \|\varphi(x - y)\|_2^2 = (1 \pm \epsilon) \|x - y\|_2^2 = (1 \pm \epsilon) \sum_{i=1}^n (x_i - y_i)^2$$

3.1 Johnson-Lindenstrauss Lemma

Lemma 3 (Distributional Johnson-Lindenstrauss 1984).

$\forall \epsilon > 0$, there is a randomized $\varphi : (R)^n \rightarrow (R)^k$ such that $\forall x, y \in (R)^n$

we have

$$P \left[\|\varphi(x) - \varphi(y)\| \in (1 \pm \epsilon) \|x - y\|_2 \right] \geq 1 - e^{-\frac{\epsilon^2 k}{9}}$$

($e^{-\frac{\epsilon^2 k}{9}}$ is the failure probability.)

In original Johnson-Lindenstrauss lemma: φ : a random k -dimensional subspace.

Proof.

Take

$$\varphi(x) = \left(\sum_{i=1}^n g_{1i} x_i, \sum_{i=1}^n g_{2i} x_i, \dots, \sum_{i=1}^n g_{ki} x_i \right) \frac{1}{\sqrt{k}}$$

Each g_{ji} is a Gaussian/normal $N(0,1)$:

$$pdf(g) = \frac{1}{2\pi} e^{-\frac{g^2}{2}}$$

Recall: What did we use to prove the correctness of Tug-of-War?

(1) $E[\sigma_i] = 0$

(2) $E[\sigma_i^2] = 1$

(3) $E[\sigma_i^4] = 1$

This is satisfied by $\sigma_i \in \{\pm 1\}$, but also by the Gaussian/normal random variable.

Consider $k = 1$:

$$\varphi(x) = \sum g_i x_i$$

Definition 4 (Stability Property).

$$\sum_{i=1}^k g_i x_i \sim \|x\|_2 \cdot a = \left(\sum x_i^2\right)^{\frac{1}{2}} \cdot a$$

a is another Gaussian $N(0,1)$

□

The probability density distribution for a centrally spherically symmetric vector $\bar{g} = (g_1, \dots, g_n)$

$$pdf(\bar{g}) = \left(\frac{1}{\sqrt{2\pi}}\right)^n \cdot e^{-\frac{g_1^2}{2}} \cdot e^{-\frac{g_2^2}{2}} \cdot \dots \cdot e^{-\frac{g_n^2}{2}} = \left(\frac{1}{\sqrt{2\pi}}\right)^n \cdot e^{-\frac{\sum_{i=1}^n g_i^2}{2}}$$

$\bar{g} \cdot x$ is distributed as $\bar{g}' \cdot (\|X\|_x, 0, 0, \dots, 0) = g'_1 \cdot \|x\|_2$

General k :

$$\|\phi(x) - \phi(y)\| = \|\phi(x - y)\| \approx \|x - y\|_2 \leftarrow \|z\|_2 \text{ where } z = x - y$$

fix z :

$$\phi(z) = \frac{1}{\sqrt{k}} \cdot \left(\sum g_{1i} z_i, \dots, \sum g_{ki} z_i\right) \sim \frac{1}{\sqrt{k}} \cdot \left(a_1 \cdot \|z\|, a_2 \cdot \|z\|, \dots, a_k \cdot \|z\|\right)$$

where each a_i is Gaussian distributed

$$\begin{aligned} \|\phi(z)\|_2^2 &= \frac{1}{k} \sum_{j=1}^k a_j^2 \cdot \|z\|^2 \\ &= \|z\|^2 \cdot \frac{1}{k} \sum_{j=1}^k a_j^2 \\ &= \|z\|^2 \cdot \mathcal{X}_k^2 \end{aligned}$$

This is \mathcal{X}^2 (Chi-squared) distributed with k degrees of freedom.

Fact:

$$P[\mathcal{X}_k^2 \notin (1 \pm \epsilon)] \leq 2 \cdot e^{-\frac{k}{4}(\epsilon^2 - \epsilon^3)}$$

for $\epsilon < \frac{1}{2}$ this gives the DJL

Corollary 5. For all N vectors $(x_1, x_2, \dots, x_N) \in \mathbb{R}^d$ in d -dimension, there exists a random ϕ from DJL such that with $k = O\left(\frac{\log(N)}{\epsilon^2}\right)$, for all $i \neq j; i, j \in [N]$:

$$P\left[\|\phi(x_i) - \phi(x_j)\| \in (1 \pm \epsilon)\|x_i - x_j\|\right] \geq 1 - \frac{1}{N}$$

Proof. Pick $k = c \cdot \frac{\log(N)}{\epsilon^2}$, DJL states:

$$\forall x, y \quad P \left[\|\phi(x) - \phi(y)\| \in (1 \pm \epsilon) \|x - y\| \right] \geq 1 - e^{-\frac{\epsilon^2 k}{9}} \geq 1 - \frac{1}{N^3}$$

by union bound:

$$P \left[\forall i, j : \|\phi(x) - \phi(y)\| \in (1 \pm \epsilon) \|x - y\| \text{ for } x = x_i, y = x_j \right] \geq 1 - \binom{N}{2} \cdot \frac{1}{N^3} \geq 1 - \frac{1}{N}$$

□

For $k \times n$ matrix \mathbb{G} and vector \mathbf{x} , where each entry in \mathbb{G} is a Gaussian:

$$\phi(x) = \frac{1}{\sqrt{k}} \cdot \mathbb{G} \cdot \mathbf{x}$$

with $1 \pm \epsilon$ approximation,

$$\phi : l_2^d \rightarrow l_2^k$$

where

$$l_2^d = \|x - y\|_2 = \sum_{j=1}^d (x_j \cdots y_j)$$

What about l_1 ?

$$l_1^d : \mathbb{R}^d \text{ where } \|x - y\|_1 = \sum_{i=1}^n |x_i - y_i|$$

$$l_p^d : \mathbb{R}^d \text{ where } \|x - y\|_1 = \left(\sum_{i=1}^n |x_i - y_i|^p \right)^{\frac{1}{p}}$$

For l_1 : N vectors into lower dimensional l_1

$$K = N^{\Omega(\frac{1}{D})} \text{ for D-approximation}$$

Alternative Sketch:

$$\phi(x) = \frac{1}{k} \cdot \mathbb{C} \cdot \mathbf{x}$$

where \mathbb{C} is a matrix with Cauchy distribution. So given $\phi(x), \phi(y)$ we can estimate $\|x - y\|$ as the median ($\|\phi(x) - \phi(y)\|$) of the absolute values of the k coordinates.

It's enough to take

$$k = O\left(\frac{\log(N)}{\epsilon^2}\right)$$

Cauchy variables are the 1-stable distribution: $\sum c_i x_i$, where c_i are random Cauchy, is distributed as $\|x\|_1 \cdot c$ where c is also Cauchy. In general, for $p \in (0, 2]$, there exist p -stable distributions satisfying the above with $\|x\|_1$ replaced by $\|x\|_p$.